IN THE CLAIMS:

The text of all pending claims, (including withdrawn claims) is set forth below. Cancelled and not entered claims are indicated with claim number and status only. The claims as listed below show added text with <u>underlining</u> and deleted text with <u>strikethrough</u>, except where strikethrough cannot be easily perceived, in which case the text of any deleted subject matter is shown by being placed within double brackets, as provided in 37 C.F.R § 1.121(c)(2). The status of each claim is indicated with one of (original), (currently amended), (cancelled), (withdrawn), (new), (previously presented), or (not entered).

1-7. (Cancelled)

8. (currently amended) An encryption device comprising XOR means and nonlinear transform means, said encryption device further comprising:

a random number generator for generating a random number h, where h is an integer between zero and q-1;

q sets of fixed \underline{mask} values $\underline{FM_{i,h}}$, where q is an integer equal to three or more, wherein equations, $FM_{0,h} = C_h XOR L1_0$ ($FMin_h$), $\underline{and} C_h = c_{h,15}c_{h,14}...c_{h,0}$, and $D_h = d_{h,15}d_{h,14}...d_{h,07}$ are satisfied, where $FM_{i,h}$ is the i-th fixed value of the h-th set of said q sets of fixed \underline{mask} values, $\underline{where} \ Fmin_h$ is a selected one of said fixed \underline{mask} values, and $\underline{where} \ i$ is an integer-and \underline{j} is an integer,

q sets of fixed <u>S-box table</u>values, wherein equations, $(c_{0,j} \text{ XOR } c_{1,j}) \lor (c_{1,j} \text{ XOR } c_{2,j}) \lor ... \lor (c_{q^-2,j} \text{ XOR } c_{q^-1,j}) = (111...11)_2$ and $(d_{0,j} \text{ XOR } d_{1,j}) \lor (d_{1,j} \text{ XOR } d_{2,j}) \lor ... \lor (d_{q^-2,j} \text{ XOR } d_{q^-1,j}) = (111...11)_2$ are satisfied, a fixed table <u>S-box table</u> before masking is defined as S[x], and i-th masked fixed table <u>S-box table</u> is defined as S_{j,h}[x] = <u>S{[x XOR c_{i,j}] XOR [[d_{i,j}]] d_{h_{i,j}}} for the j-th fixed value, where j is an integer;</u>

linear transform means $L1_i(x)$ and linear transform means $L2_i(x)$, wherein the linear transform means $L1_i(x)$, the nonlinear transform means with the masked fixed table S_{j,h} [x] and the linear transform means $L2_i(x)$ operate in i-th one of rounds; and

a first selector for selecting one fixed value of the h-th set of said q sets of fixed <u>mask</u> values in response to the random number h,

said XOR means XORing an input thereto with an XOR of a key with said selected fixed value.

9-34. (Cancelled)

35. (Previously presented) The encryption device according to claim 8, characterized in that the linear transform means $L1_i(x)$ is defined as $L1_i(x) = x$.

- 36. (Previously presented) The encryption device according to claim 8, characterized in that the linear transform means $L2_i(x)$ is defined as $L2_i(x) = MixedColumn(Shift(x))$.
- 37. (Previously presented) The encryption device according to claim 8, characterized in that the linear transform means $L2_i(x)$ is defined as Shift(x).
- 38. (currently amended) An encryption device comprising XOR means and nonlinear transform means, said encryption device further comprising:

a random number generator for generating a random number h, where h is an integer between zero and q-1;

q sets of fixed <u>mask_values_FM_{i,h}</u>, where q is an integer <u>equal to three</u> or more, wherein equations, FM_{i,h} = C_h XOR L1i(L2i-1(D_h)) for i≥1, C_h = c_{h,15}c_{h,14}...c_{h,0}, and D_h = d_{h,15}d_{h,14}...d_{h,0}, are satisfied, where i is an integer-and j is an integer,

q sets of fixed <u>S-box table-values</u>, where equations, $(c_{0,j} \text{ XOR } c_{1,j}) \lor (c_{1,j} \text{ XOR } c_{2,j}) \lor ... \lor (c_{q-2,j} \text{ XOR } c_{q-1,j}) = (1111...11)_2$ and $(d_{0,j} \text{ XOR } d_{1,j}) \lor (d_{1,j} \text{ XOR } d_{2,j}) \lor ... \lor (d_{q-2,j} \text{ XOR } d_{q-1,j}) = (1111...11)_2$, are satisfied, and a fixed table <u>S-box table</u> before masking is defined as S[x] and an i-th masked fixed table <u>S-box table</u> is defined as $S_{j,h}[x] = [x \text{ XOR } c_{i,j}] \text{ XOR } d_{i,j}$ for the j-th fixed value, where j is an integer;

linear transform means $L1_i(x)$ and linear transform means $L2_i(x)$, wherein the linear transform means $L1_i(x)$, the nonlinear transform means with the masked fixed table S-box table $S_{i,h}[x]$ and the linear transform means $L2_i(x)$ operate in i-th one of rounds; and

a first selector for selecting one fixed value of the h-th set of said q sets of fixed mask values in response to the random number h;

said XOR means XORing an input thereto with an XOR of a key with said selected fixed value.

39. (Previously presented) The encryption device according to claim 38, characterized in that the linear transform means $L1_i(x)$ is defined as $L1_i(x) = x$.

- 40. (Previously presented) The encryption device according to claim 38, characterized in that the linear transform means $L2_i(x)$ is defined as $L2_i(x) = MixedColumn(Shift(x))$.
- 41. (Previously presented) The encryption device according to claim 38, characterized in that the linear transform means $L2_i(x)$ is defined as Shift(x).

42-49. (Cancelled)

50. (currently amended) An encryption device comprising XOR means and nonlinear transform means, said encryption device further comprising:

a random number generator for generating a random number h, where h is an integer between zero and q-1;

q sets of masked fixed \underline{mask} values $\underline{FM_{i,h}}$ and \underline{q} sets of fixed tables, where \underline{q} is an integer equal to three or more, wherein equations, $FM_{0,h} = C_h XOR L1_0$ ($FMin_h$), $\underline{and} C_h = c_{h,15}c_{h,14}...c_{h,0}$, and $D_h = d_{h,15}d_{h,14}...d_{h,07}$ are satisfied, where $FM_{i,h}$ is the i-th fixed value of the h-th set of said \underline{q} sets of fixed \underline{mask} values, $\underline{where} FMin_h$ is a selected one of said fixed \underline{mask} values, and \underline{w} where i is an integer and \underline{j} is an integer,

q sets of fixed <u>S-box table-values</u>, wherein equations[[,]] ($c_{0,j}$ XOR $c_{1,j}$) \vee ($c_{1,j}$ XOR $c_{2,j}$) $\vee ... \vee$ ($c_{q^-2,j}$ XOR $c_{q^-1,j}$) = (111...11)₂ and ($d_{0,j}$ XOR $d_{1,j}$) \vee ($d_{1,j}$ XOR $d_{2,j}$) $\vee ... \vee$ ($d_{q^-2,j}$ XOR $d_{q^-1,j}$) = (111...11)₂ are satisfied, a fixed table <u>S-box table</u> before masking is defined as S[x], and i-th masked fixed table <u>S-box table</u> is defined as S_{j,h}[x] = [x XOR c[[i]]_{h,j}] XOR d[[i]]_{h,j} for the j-th fixed value, where j is an integer;

a selector for selecting one of said q sets of fixed table S-box tables in response to the random number h,

said nonlinear transform means nonlinearly transforming an input thereto in accordance with said selected set of fixed tableS-box tables; and

a plurality of encrypting rounds, wherein i-th one of said plurality of encrypting rounds comprises the XOR means, the fixed tableS-box tables, the selector, linear transform means $L1_i(x)$ and linear transform means $L2_i(x)$, for that round, and wherein the fixed tableS-box tables for said plurality of respective encrypting rounds are identical, and wherein the linear transform means $L1_i(x)$, the nonlinear transform means with the masked fixed tableS-box table $S_{i,h}[x]$ and the linear transform means $L2_i(x)$ operate in that round.

- 51. (Previously presented) The encryption device according to claim 50, characterized in that the linear transform means $L1_i(x)$ is defined as $L1_i(x) = x$.
- 52. (Previously presented) The encryption device according to claim 50, characterized in that the linear transform means $L2_i(x)$ is defined as $L2_i(x) = MixedColumn(Shift(x))$.
- 53. (Previously presented) The encryption device according to claim 50, characterized in that the linear transform means $L2_i(x)$ is defined as Shift(x).
- 54. (currently amended) An encryption device comprising XOR means and nonlinear transform means, said encryption device further comprising:

a random number generator for generating a random number h, where h is an integer between zero and q-1;

q sets of masked fixed <u>mask</u> values <u>FM_{i,h}</u> and q sets of fixed tables, where q is an integer equal to three or more, wherein equations, $FM_{i,h} = C_h XOR L[[1]]2_i(L[[2]]1_{i-1}(D_h))$ for $i \ge 1$, $C_h = c$ $h_{1,15}C_{h,14}...c_{h,0}$, and $D_h = d_{h,15}d_{h,14}...d_{h,0}$, are satisfied, where $FM_{i,h}$ is the i-th fixed value of the h-th set of said q sets of fixed <u>mask</u> values, where i is an integer-and j is an integer,

q sets of fixed <u>S-box table</u>values, wherein equations, $(c_{0,j} \text{ XOR } c_{1,j}) \lor (c_{1,j} \text{ XOR } c_{2,j}) \lor ... \lor (c_{q^-2,j} \text{ XOR } c_{q^-1,j}) = (111...11)_2$ and $(d_{0,j} \text{ XOR } d_{1,j}) \lor (d_{1,j} \text{ XOR } d_{2,j}) \lor ... \lor (d_{q^-2,j} \text{ XOR } d_{q^-1,j}) = (111...11)_2$ are satisfied, a fixed table <u>S-box table</u> before masking is defined as S[x], and i-th masked fixed table <u>S-box table</u> is defined as $S_{j,h}[x] = [x \text{ XOR } c[[i]]_{h,j}]$ XOR $d[[i]]_{h,j}$ for the j-th fixed value, where j is an integer;

a selector for selecting one of said q sets of fixed tableS-box tables in response to the random number h,

said nonlinear transform means nonlinearly transforming an input thereto in accordance with said selected set of fixed tableS-box tables; and

a plurality of encrypting rounds, wherein i-th one of said plurality of encrypting rounds comprises the XOR means, the fixed tableS-box tables, the selector, linear transform means $L1_i(x)$ and linear transform means $L2_i(x)$, for that round, and wherein the fixed tableS-box tables for said plurality of respective encrypting rounds are identical, and wherein the linear transform means $L1_i(x)$, the nonlinear transform means with the masked fixed tableS-box table $S_j[x]$ and the linear transform means $L2_i(x)$ operate in that round.

- 55. (Previously presented) The encryption device according to claim 54, characterized in that the linear transform means $L1_i(x)$ is defined as $L1_i(x) = x$.
- 56. (Previously presented) The encryption device according to claim 54, characterized in that the linear transform means $L2_i(x)$ is defined as $L2_i(x) = MixedColumn(Shift(x))$.
- 57. (Previously presented) The encryption device according to claim 54, characterized in that the linear transform means $L2_i(x)$ is defined as Shift(x).

58-63. (Cancelled)

64. (currently amended) An encryption device comprising a random number generator for generating a random number h, where h is an integer between zero and q-1, and a first plurality of encrypting rounds, wherein

i-th one of said plurality of encrypting rounds comprises nonlinear transform means for nonlinearly transforming an input thereto, and XOR means for XORing a first input thereto with a second input thereto for that round, where i is an integer;

the second input to said XOR means is coupled to an output of said nonlinear transform means; and

said nonlinear transform means comprises:

q sets of fixed $\underline{\mathsf{mask}}$ values $\underline{\mathsf{FM}_{i,h}}$, where q is an integer equal to three or more, wherein equations, $\mathsf{FM}_{0,h} = \mathsf{C}_h \; \mathsf{XOR} \; \mathsf{L1}_0 \; (\mathsf{FMin}_{\underline{h}})$, $\underline{\mathsf{and}} \; \mathsf{C}_h = \mathsf{c}_{h,15} \mathsf{c}_{h,14} \dots \mathsf{c}_{h,0}$, $\underline{\mathsf{and}} \; \mathsf{D}_{\underline{h}} = \mathsf{d}_{h,15} \mathsf{d}_{h,14} \dots \mathsf{d}_{h,0}$, are satisfied, where $\mathsf{FM}_{i,h}$ is the i-th fixed value of the h-th set of said q sets of fixed $\underline{\mathsf{mask}}$ values, where $\underline{\mathsf{FMin}_h}$ is a selected one of said fixed $\underline{\mathsf{mask}}$ values, and $\underline{\mathsf{j}}$ is an integer,

q sets of fixed <u>S-box table-values</u>, wherein equations, $(c_{0,j} \text{ XOR } c_{1,j}) \lor (c_{1,j} \text{ XOR } c_{2,j}) \lor ... \lor (c_{q^-2,j} \text{ XOR } c_{q^-1,j}) = (111...11)_2$ and $(d_{0,j} \text{ XOR } d_{1,j}) \lor (d_{1,j} \text{ XOR } d_{2,j}) \lor ... \lor (d_{q^-2,j} \text{ XOR } d_{q^-1,j}) = (111...11)_2$ are satisfied, a fixed table <u>S-box table</u> before masking is defined as S[x], and i-th masked fixed table <u>S-box table</u> is defined as $S_{j,h}[x] = [x \text{ XOR } c[[j]]_{h,j}] \text{ XOR } d[[j]]_{h,j}$ for the j-th fixed value, where j is an integer;

a selector for selecting one of said q sets of fixed <u>mask</u> values in response to the random number h;

further XOR means for XORing an input thereto with an XOR of a key with said selected fixed value:

linear transform means $L1_i(x)$;

a plurality of nonlinear transform means for nonlinearly transforming an input in accordance with a fixed table S-box table;

linear transform means L2_i(x), wherein the linear transform means L1_i(x), the nonlinear transform means with the masked fixed table S-box table S_{j,h} [x] and the linear transform means L2_i(x) operate in that round; and

a selector for selecting one of said plurality of nonlinear transform means.

- 65. (Previously presented) The encryption device according to claim 64, characterized in that the linear transform means $L1_i(x)$ is defined as $L1_i(x) = x$.
- 66. (Previously presented) The encryption device according to claim 64, characterized in that the linear transform means $L2_i(x)$ is defined as $L2_i(x) = MixedColumn(Shift(x))$.
- 67. (Previously presented) The encryption device according to claim 64, characterized in that the linear transform means $L2_i(x)$ is defined as Shift(x).
- 68. (currently amended) An encryption device comprising a random number generator for generating a random number h, where h is an integer between zero and q-1, and a first plurality of encrypting rounds, wherein

i-th one of said plurality of encrypting rounds comprises nonlinear transform means for nonlinearly transforming an input thereto, and XOR means for XORing a first input thereto with a second input thereto for that round, where i is an integer;

the second input to said XOR means is coupled to an output of said nonlinear transform means; and

said nonlinear transform means comprises:

q sets of fixed $\underline{\text{mask}}$ values $\underline{\text{FM}_{i,h}}$, where q is an integer equal to three or more, wherein equations[[,]] $\underline{\text{FM}_{i,h}} = C_h \, \text{XOR L1i}(L2i\text{-}1(D_h))$ for $i \ge 1$, $C_h = c_{h,15} c_{h,14} \dots c_{h,0}$, and $D_h = d_{h,15} d_{h,14} \dots d_{h,0}$, are satisfied, where $\underline{\text{FM}_{i,h}}$ is the i-th fixed value of the h-th set of said q sets of fixed $\underline{\text{mask}}$ values, where j is an integer,

q sets of fixed <u>S-box table-values</u>, wherein equations[[,]] ($c_{0,j}$ XOR $c_{1,j}$) \vee ($c_{1,j}$ XOR $c_{2,j}$) $\vee ... \vee$ ($c_{q^-2,j}$ XOR $c_{q^-1,j}$) = (111...11)₂ and ($d_{0,j}$ XOR $d_{1,j}$) \vee ($d_{1,j}$ XOR $d_{2,j}$) $\vee ... \vee$ ($d_{q^-2,j}$ XOR $d_{q^-1,j}$) = (111...11)₂ are satisfied, a fixed table <u>S-box table</u> before masking is defined as S[x], and i-th

masked fixed table S-box table is defined as $S_{j,h}[x] = [x] \times ([i])_{h,j} \times ([i])_{h,j}$ for the j-th fixed value, where j is an integer;

a selector for selecting one of said q sets of fixed <u>mask</u> values in response to the random number h;

further XOR means for XORing an input thereto with an XOR of a key with said selected fixed value;

linear transform means $L1_i(x)$;

a plurality of nonlinear transform means for nonlinearly transforming an input in accordance with a fixed tableS-box table;

linear transform means $L2_i(x)$, wherein the linear transform means $L1_i(x)$, the nonlinear transform means with the masked fixed table S-box table $S_{j,h}[x]$ and the linear transform means $L2_i(x)$ operate in that round; and

a selector for selecting one of said plurality of nonlinear transform means.

- 69. (Previously presented) The encryption device according to claim 68, characterized in that the linear transform means $L1_i(x)$ is defined as $L1_i(x) = x$.
- 70. (Previously presented) The encryption device according to claim 68, characterized in that the linear transform means $L2_i(x)$ is defined as $L2_i(x) = MixedColumn(Shift(x))$.
- 71. (Previously presented) The encryption device according to claim 68, characterized in that the linear transform means $L2_i(x)$ is defined as Shift(x).
- 72. (currently amended) A program stored on a computer or machine-readable storage medium for use in an encryption device, said program operable to effect the steps of: selecting one set of q sets of fixed mask-values_FM_i,h, where q is an integer equal to three or more, in response to a random number h, where h is an integer between zero and q-1;

XORing an input value with an XOR of a key with said selected fixed value in i-th one of rounds, wherein equations, $FM_{0,h} = C_h$ XOR L1₀ ($FMin_h$)₇ and $C_h = c_{h,15}c_{h,14}...c_{h,0}$, and $D_h = d_{h,15}d$ $h,14...d_{h,07}$ are satisfied, where $FM_{i,h}$ is the i-th fixed value of the h-th set of said q sets of fixed mask values, where $FM_{i,h}$ is a selected one of said fixed mask values and where i is an integer and j is an integer;

selecting one set $S_{j,h}$ [x] of q sets of masked fixed table S-box tables in response to the

random number h in that round, wherein equations, $(c_{0,j} \text{ XOR } c_{1,j}) \lor (c_{1,j} \text{ XOR } c_{2,j}) \lor ... \lor (c_{q^-2,j} \text{ XOR } c_{q^-1,j}) = (111...11)_2$ and $(d_{0,j} \text{ XOR } d_{1,j}) \lor (d_{1,j} \text{ XOR } d_{2,j}) \lor ... \lor (d_{q^-2,j} \text{ XOR } d_{q^-1,j}) = (111...11)_2$ are satisfied, a fixed table before masking is defined as S[x], and i-th masked fixed table so table is defined as $S_{j,h}[x] = [x] \times SOR$ c[[i]]_{h,j} XOR d[[i]]_{h,j} for the j-th fixed value, where j is an integer;

nonlinearly transforming an input value in accordance with said selected set $S_{j,h}$ [x] of fixed tableS-box tables in that round; and

linearly transforming an input value and the nonlinearly transformed value in that round.